

I) What is our research problem? The **satisfiability problem** is one of the most fundamental problems studied in computer science, with multiple applications in knowledge representation [17], formal verification [20], and planning [11]. It has also been a central undertaking of mathematical logic since the publication of *Hilbert's Grundzüge der theoretischen Logik*. This problem is parametrised by a formal language \mathcal{L} (a.k.a. *logic*, e.g. first-order logic), and asks if an input formula φ (written in the syntax of \mathcal{L}) is satisfied in *some* structure, i.e. it has a *model*. Formally:

(Finite) Satisfiability Problem	(Finite) Entailment Problem
<i>Parameters:</i> A logic \mathcal{L} .	Logics \mathcal{L}_1 and \mathcal{L}_2 .
<i>Input:</i> An \mathcal{L} -formula φ .	An \mathcal{L}_1 -formula φ_1 and an \mathcal{L}_2 -formula φ_2 .
<i>Question:</i> Does φ have a (finite) model?	Do all (finite) models of φ_1 satisfy φ_2 ?

The satisfiability problem generalises to the above-defined entailment problem.¹ In the restricted case of *finite* satisfiability and entailment problems, we are interested in models of finite sizes only. These problems have served as a yardstick for measuring the “quality” of logics, driving the discovery of decidable first-order fragments [5, 24] in the last decades. Satisfiability underlies symbolic program verification [1], while (query) entailment is the main reasoning problem over incomplete databases in Ontology-Based Data Access (OBDA) [23]. In formal verification, the input formula expresses a system specification, and satisfiability determines its feasibility. In OBDA, we seek certain (true in all extensions) query answers over incomplete data. There, φ_1 encodes a database \mathcal{D} together with some integrity constraints, while φ_2 represents the query. Entailment checks whether all “valid” completions of \mathcal{D} (w.r.t. φ_1) satisfy the query φ_2 .

II) On what logics do we focus? FINGO focuses on families of *dynamic logics* and query languages supporting navigational features, and studies their finite satisfiability and entailment problems. In the case of decidability, we also aim to establish tight complexity bounds for these problems.

We study logics supporting regular path expressions, focusing on (extensions of) well-known program verification logics like Propositional Dynamic Logic (PDL) and the μ -calculus.

Regular path expressions have become the de facto standard for querying graph-structured data, with languages like Cypher, Gremlin, and PGQL widely adopted in industry [2]. They also arise naturally in bioinformatics and cheminformatics [10, 21], where domain experts frequently seek associations among entities in protein, cellular, drug, and disease networks (stored as graphs). In this setting, identifying gene–disease–drug associations (paths in the data) can aid in the discovery of new treatment strategies. Consult the full proposal for more motivation behind the choice of PDL [9] and the μ -calculus [19] as the main object of study of FINGO.

III) What are the existing approaches? Dynamic logics have the *tree-like (counter)model property*, i.e. whenever φ_1 does not entail φ_2 , there exists an infinite “tree-representable” model of φ_1 violating φ_2 . Hence, one can benefit from the theory of formal languages and solve the unrestricted satisfiability and entailment problems via **automata-based methods**. One relies on Alternating Parity Tree Automata (APTA) [25], finite-state machines that read infinite node-labelled trees top-down (accepting or rejecting a tree). An algorithm solving the entailment problem via APTA is given on the next page. Despite its simplicity, its adaptation to new settings presents two key challenges. First, defining the right notion of automata and solving its non-emptiness problem.

¹Note that φ entails ψ if and only if $\varphi \wedge \neg\psi$ is unsatisfiable.

Second, designing a suitable encoding of models as tree-like structures.

Algorithm 1: Deciding the entailment problem for the input formulæ φ_1 and φ_2 .

- 1 Construct APTA \mathcal{A}_1 and \mathcal{A}_2 accepting all tree-like models of φ_1 and φ_2 .
 - 2 Construct an APTA \mathcal{A}_2^- accepting all tree-like structures violating φ_2 . // The APTA \mathcal{A}_2^- is usually constructed through projection and complementation of \mathcal{A}_2 .
 - 3 Return YES if the *intersection* of the languages \mathcal{A}_1 and \mathcal{A}_2^- is empty, and NO otherwise.
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For the satisfiability problem, we build \mathcal{A}_1 and test for its non-emptiness. Algorithm 1 scales well with generalisations of APTA, and was applied by M. Ortiz (the PI's mentor) to solve the entailment problem for Conjunctive Regular Path Queries (CRPQs, an expressive navigational query language) and expressive extensions of PDL [6, 22]. We stress that the correctness of this method heavily relies on the existence of infinite tree-like (counter)models encodable via APTA.

IV) What about finite-model reasoning? Unfortunately, the presented technique **fails** when moving to the finite-model reasoning case, as finite models for formulæ written in (extensions of) PDL and the μ -calculus are unlikely to be tree-like. As an example, abstractions of computer programs as well as real-world systems in biochemistry are inherently finite and frequently involve complex motifs like cycles and cliques, absent in tree-like structures. Despite the broad applicability of dynamic logics, only tiny progress has been made in the past on their finite satisfiability and entailment problems. We believe this is largely due to the absence of available automata-based methods,² crucial for dealing with regular expressions or fixed-points. We stress:

Due to the lack of automata-based tools for finite-model reasoning, little progress has been made on finite satisfiability and entailment problems for dynamic logics over the past 30 years.

The above claim refers to logics for which finite and unrestricted satisfiability or entailment do not coincide and recent work by Figueira [8]). The *only known exceptions* are the entailment of CRPQs over \mathcal{ALC} (PDL without regular expressions), a recent breakthrough by Gutowski et al. [15, 16] established after a series of mathematically laborious works [12, 13, 14], and the finite satisfiability problem for the μ -calculus with converse by M. Bojańczyk [3]. The second paper is based on a simple yet ingenious idea: Bojańczyk replaced APTA by a dedicated finite-graph operating automaton and proved decidability of the non-emptiness problem. His approach effectively solves the finite satisfiability for the logic, but unfortunately does not extend to richer logics, like the graded μ -calculus with converse [4] (for which the finite satisfiability problem has been open for over two decades), or to the entailment problem (*e.g.* of CRPQs). This is where FINGO aims to make a unique and significant contribution.

V) What are the main goals of FINGO? The main goal of FINGO can be summarised as follows.

The mission of FINGO is to **alter the landscape of available techniques** for finite-model reasoning and finite entailment problems, and apply them to **solve long-standing open problems** in the area of dynamic logics, with an emphasis on (extensions of) Propositional Dynamic Logic and the μ -calculus. FINGO will achieve this goal by **designing and providing complexity results for novel automata models operating on finite graphs**, extending APTA.

²Other possible tools are tableaux methods and cyclic proof systems, but they can be reformulated as automata.

The main research directions of FINGO are structured into work packages (WPs) outlined below. Each work package develops algorithms for novel finite-graph automata, motivated by their applications in OBDA and decision procedures for dynamic logic.

(WP1) Adapt Algorithm 1 to the setting of finite-graph automata and OBDA.

(WP2) Solve the non-emptiness problem for two-way finite-graph automata with counting.

(WP3) Understand the limits of automata-based approaches to the finite satisfiability problem by analysing extensions of finite-graph automata with pebbles, counters, or storage.

(WP1) will provide a uniform algorithm for navigational query answering in OBDA under the finite model semantics. (WP2) will establish a novel automata-based technique for finite-model reasoning for logics with inverses, counting, and regular expressions, and will settle the long-standing open problem of finite satisfiability for the two-way graded μ -calculus. Finally, (WP3) will provide partial progress on techniques for generalisations of PDL like LoopPDL [7] or Non-Regular PDL [18], whose finite satisfiability problems have remained open for over 40 years.

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